

# The moment-SOS approach II: Some Applications

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# A preliminary result

- Let  $\mathbf{B} \subset \mathbb{R}^n$  be a **simple set** like e.g., a Box  $[-1, 1]^n$  or an ellipsoid,
- $\mathbf{K} := \{ (x, y) : x \in \mathbf{B}; g_j(x, y) \geq 0, j = 1, \dots, m \}$  and  $h : \mathbb{R}^n \times \mathbb{R}^p \rightarrow \mathbb{R}$ .

Goal: Approximate the function  $F : \mathbf{B} \rightarrow \mathbb{R}$

$$x \mapsto F(x) := \sup_y \{ h(x, y) : (x, y) \in \mathbf{K} \}, \quad x \in \mathbf{B}.$$

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... by **POLYNOMIALS**  $J_k \in \mathbb{R}[x]$  of increasing degree ... and with some **guarantee of convergence**.

Fix  $k \in \mathbb{N}$  and consider the optimization problem:

$$\inf_{J \in \mathbb{R}[x]_k} \left\{ \int_{\mathbf{B}} (J(x) - F(x)) dx : J(x) \geq F(x), \quad \forall x \in \mathbf{B} \right\}$$

Or, equivalently:

$$\inf_{J \in \mathbb{R}[x]_k} \{ \|J - F\|_1 : p(x) \geq F(x), \quad \forall x \in \mathbf{B} \}$$

where  $\|\cdot\|_1$  is the  $L_1$ -norm

$$h \mapsto \|h\|_1 := \int_{\mathbf{B}} |h(x)| dx, \quad h \in L_1(\mathbf{B}).$$

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To fix ideas, let  $\mathbf{B} := [-1, 1]^n = \{x : x_i^2 \leq 1, i = 1, \dots, n\}$  and set  $g_0(x) = 1$  and  $g_{m+i}(x) = 1 - x_i^2, i = 1, \dots, n$ .

### A simple strategy:

Replace the difficult constraint  $J(x) - F(x), \forall x \in \mathbf{B}$  with the SOS-based positivity certificate

$$J(x) - h(x, y) = \sum_{j=0}^{m+n} \sigma_j(x, y) g_j(x, y),$$

for some SOS polynomials  $(\sigma_j)$  such that  $\deg(\sigma_j g_j) \leq 2k$ .

Indeed

$$J(x) \geq h(x, y), \quad \forall (x, y) \in \mathbf{K}; \quad x \in \mathbf{B} \Rightarrow J(x) \geq F(x), \quad \forall x \in \mathbf{B}.$$

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Then solve:

$$\inf_{J, \sigma_j} \left\{ \int_{\mathbf{B}} (J(x) - F(x)) dx : \right.$$

$$\text{s.t. } J(x) - h(x, y) = \sum_{j=0}^{m+n} \sigma_j(x, y) g_j(x, y), \quad \forall (x, y) \in \mathbb{R}^{n+p}$$

$$J \in \mathbb{R}[\mathbf{x}]_k; \sigma_j \text{ SOS; } \deg(\sigma_j g_j) \leq 2k \}.$$

Theorem

$\Rightarrow$  This is an SDP and it has an optimal solution  $J_k^* \in \mathbb{R}[x]_k$ .  
Moreover,  $\|J_k^* - f\|_1 \rightarrow 0$  as  $k \rightarrow \infty$ .




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Write  $J(x) = \sum_{\alpha \in \mathbb{N}^n} J_\alpha x^\alpha$ .

I.  The criterion  $\int_{\mathbf{B}} (J(x) - f(x)) dx$  reads:

$$\sum_{\alpha \in \mathbb{N}^n} J_\alpha \underbrace{\int_{\mathbf{B}} x^\alpha dx}_{=\gamma_\alpha \text{ known}} - \underbrace{\int_{\mathbf{B}} F(\mathbf{x}) dx}_{\text{an (unknown) constant } c} = \sum_{\alpha \in \mathbb{N}^n} J_\alpha \gamma_\alpha - c,$$

is **LINEAR** in the unknown coefficients  $J_\alpha$  of  $J \in \mathbb{R}[x]_k$ .

## II. 🖱️ The constraint

$$J(x) - h(x, y) = \sum_{j=0}^{m+n} \sigma_j(x, y) g_j(x, y),$$

reduces to

- **LINEAR** constraints on the coefficients  $J_\alpha$  and  $\sigma_{j\alpha}$ ,
- + semidefinite constraints to state that the  $\sigma_j$ 's are SOS.

# I. Approximation of sets with quantifiers

Let  $f \in \mathbb{R}[x, y]$  and let  $\mathbf{K} \subset \mathbb{R}^n \times \mathbb{R}^p$  be the semi-algebraic set:

$$\mathbf{K} := \{(x, y) : g_j(x, y) \geq 0, \quad j = 1, \dots, m\},$$

and let  $\mathbf{B} \subset \mathbb{R}^n$  be the Euclidean unit ball or the box  $[-1, 1]^n$ .

Suppose that one wants to approximate the set:

$$R_f := \{x \in \mathbf{B} : f(x, y) \leq 0 \text{ for all } y \text{ such that } (x, y) \in \mathbf{K}\}$$

as closely as desired by a sequence of sets of the form:

$$\Theta_k := \{x \in \mathbf{B} : J_k(x) \leq 0\}$$

for some polynomials  $J_k$ .

Recall that  $\mathbf{B} = \{x : 1 - x_i^2 \geq 0\}$  and set  $g_{m+i}(x) := 1 - x_i^2$ ,  $i = 1, \dots, n$ . With  $g_0 = 1$  and with  $\mathbf{K} \subset \mathbb{R}^n \times \mathbb{R}^p$  and  $k \in \mathbb{N}$ , let

$$Q_k(g) = \left\{ \sum_{j=0}^{m+n} \sigma_j(x, y) g_j(x, y) : \sigma_j \in \Sigma[x, y], \deg(\sigma_j g_j) \leq 2k \right\}$$

Let  $x \mapsto F(x) := \max \{f(x, y) : (x, y) \in \mathbf{K}\}$ , and

for every integer  $k$  consider the optimization problem:

$$\rho_k = \min_{J \in \mathbb{R}[x]_k} \left\{ \int_{\mathbf{B}} (J - F) dx : J(x) - f(x, y) \in Q_k(g); x \in \mathbf{B} \right\}$$

Remember that  $J(x) - f(x, y) \in Q_k(g)$  implies:

$$J(x) \geq f(x, y), \quad \forall (x, y) \in \mathbf{K} \quad \Rightarrow \quad J(x) \geq F(x) \quad \forall x \in \mathbf{B}.$$

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## 1. The criterion

$$\int_{\mathbf{B}} (J - F) dx = \underbrace{\int_{\mathbf{B}} -F dx}_{\text{unknown but constant}} + \sum_{\alpha} J_{\alpha} \underbrace{\int_{\mathbf{B}} \mathbf{x}^{\alpha} dx}_{\text{easy to compute}}$$

is **LINEAR** in the coefficients  $J_{\alpha}$  of the unknown polynomial  
 $J \in \mathbb{R}[\mathbf{x}]_k!$

## 2. The constraint

$$J(x) - f(x, y) = \sum_{j=0}^m \sigma_j(x, y) g_j(x, y)$$

is just **LINEAR CONSTRAINTS + LMIs!**



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Hence, the optimization problem

$$\rho_k = \min_{J \in \mathbb{R}[x]_k} \left\{ \int_{\mathbf{B}} (J - F) dx : J(\mathbf{x}) - f(\mathbf{x}, y) \in Q_k(g) \right\}$$

IS AN SDP! Moreover, it has an optimal solution  $J_k^* \in \mathbb{R}[x]_k$ !

- Alternatively, if one uses LP-based positivity certificates for  $J(\mathbf{x}) - f(\mathbf{x}, y)$ , one ends up with solving an LP!

From the definition of  $J_k^*$ , the sublevel sets

$$\Theta_k := \{x \in \mathbf{B} : J_k^*(x) \leq 0\} \subset R_f, \quad k \in \mathbb{N},$$

provide a nested sequence of INNER approximations of  $R_f$ .

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## Theorem (Lass)

(Strong) convergence in  $L_1(\mathbf{B})$ -norm takes place, that is:

$$\lim_{k \rightarrow \infty} \int_{\mathbf{B}} |J_k^* - F| dx = 0$$

and, if in addition the set  $\{x \in \mathbf{B} : F(x) = 0\}$  has Lebesgue measure zero, then

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# Ex: Polynomial Matrix Inequalities: (with D. Henrion)

Let  $x \mapsto \mathbf{A}(x) \in \mathbb{R}^{p \times p}$  where  $\mathbf{A}(x)$  is the **matrix-polynomial**

$$x \mapsto \mathbf{A}(x) = \sum_{\alpha \in \mathbb{N}^n} \mathbf{A}_\alpha x^\alpha \quad \left( = \sum_{\alpha \in \mathbb{N}^n} \mathbf{A}_\alpha x_1^{\alpha_1} \cdots x_n^{\alpha_n} \right).$$

for finitely many **real symmetric matrices**  $(\mathbf{A}_\alpha)$ ,  $\alpha \in \mathbb{N}^n$ .

... and suppose one wants to approximate the set

$$R_{\mathbf{A}} := \{x \in \mathbf{B} : \mathbf{A}(x) \succeq 0\} = \{x : \lambda_{\min}(\mathbf{A}(x)) \geq 0\}.$$

Then:

$$R_{\mathbf{A}} = \left\{ x \in \mathbf{B} : \underbrace{y^T \mathbf{A}(x) y}_{f(x,y)} \geq 0, \quad \forall y \text{ s.t. } \|y\|^2 = 1 \right\}$$

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# Illustrative example (continued)

Let  $\mathbf{B}$  be the unit disk  $\{\mathbf{x} : \|\mathbf{x}\| \leq 1\}$  and let:

$$R_{\mathbf{A}} := \left\{ \mathbf{x} \in \mathbf{B} : \mathbf{A}(\mathbf{x}) \left( = \begin{bmatrix} 1 - 16x_1x_2 & x_1 \\ x_1 & 1 - x_1^2 - x_2^2 \end{bmatrix} \right) \succeq 0 \right\}$$

Then by solving relatively simple **semidefinite programs**, one may approximate  $R_{\mathbf{A}}$  with **sublevel sets** of the form:

$$\Theta_k := \{x \in \mathbf{B} : J_k^*(x) \geq 0\}$$

for some polynomial  $J_k^*$  of degree  $k = 2, 4, \dots$  and with

$$\text{VOL}(R_{\mathbf{A}} \setminus \Theta_k) \rightarrow 0 \quad \text{as } k \rightarrow \infty.$$

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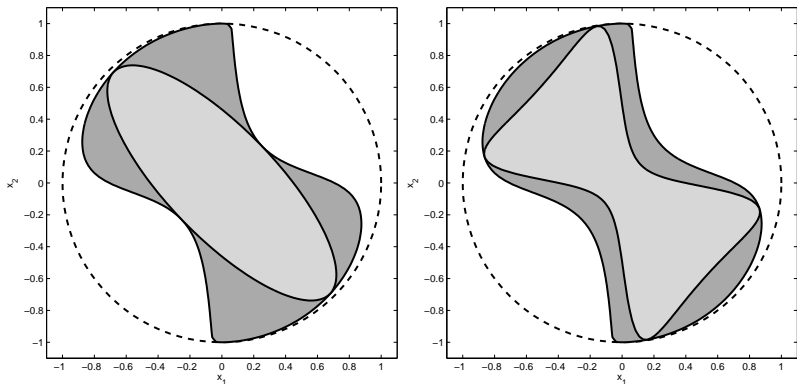
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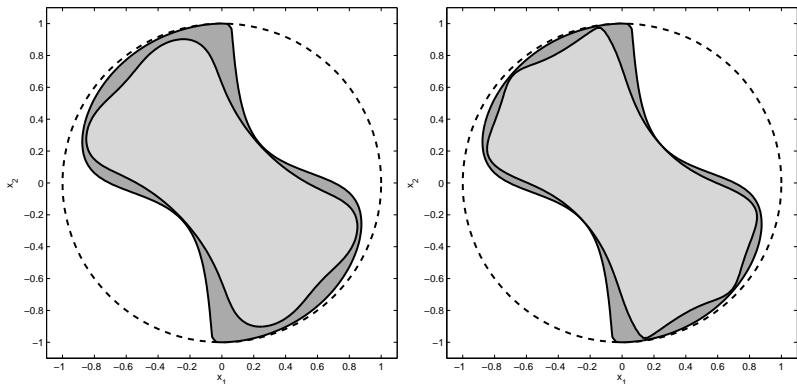
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$\Theta_2$  (left) and  $\Theta_4$  (right) inner approximations (light gray) of (dark gray) embedded in unit disk  $\mathbf{B}$  (dashed).



$\Theta_6$  (left) and  $\Theta_8$  (right) inner approximations (light gray) of (dark gray) embedded in unit disk  $\mathbf{B}$  (dashed).

## II. Convex Underestimators of Polynomials

A typical **MINLP** is of the form:

$$\mathbf{P} : \quad \inf_{\mathbf{x}} \{ f(\mathbf{x}) : \mathbf{x} \in \mathbf{K} \cap \mathbf{B}; x_i \in \{0, 1\}, \forall i \in I \},$$

where

- $\mathbf{K} := \{ \mathbf{x} : g_j(\mathbf{x}) \leq 0, j = 1, \dots, m \}$  for some  $(g_j) \subset \mathbb{R}[\mathbf{x}]$ ,
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- In the context of **large scale MINLP** the most efficient & popular strategy is to use **BRANCH & BOUND** combined with efficient **LOWER BOUNDING** techniques used at each node of the search tree.
- Typically,  $f$  is a sum  $\sum_k f_k$  where each  $f_k$  “sees” only very few variables (say 3, 4). The same observation is true for each  $g_j$  in the constraints:

Hence a very appealing idea is to pre-compute **CONVEX UNDERESTIMATORS**  $\hat{f}_k \leq f_k$  and  $\hat{g}_j \leq g_j$  for each non convex  $f_k$  and each non convex  $g_j$ , independently and separately!

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- Then at each node of the search tree of the B & B tree one computes a **LOWER BOUND** by solving convex optimization problems the form:

$$\hat{\mathbf{P}} : \inf_{\mathbf{x}} \left\{ \sum_k \hat{f}_k(\mathbf{x}) : \hat{g}_j(\mathbf{x}) \leq 0, j = 1, \dots, m; \mathbf{x} \in \mathbf{B} \right\}$$

where some of the integer variables are fixed at 0 or 1.

and then one explores the search tree.

Hence one has to solve the generic problem

Compute a "tight" convex polynomial underestimator  $p \leq f$  of a non convex polynomial  $f$  on a box  $\mathbf{B} \subset \mathbb{R}^n$ .

Message:

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## I: Characterizing convex polynomial underestimators

①  $p(\mathbf{x}) \leq f(\mathbf{x})$  for every  $\mathbf{x} \in \mathbf{B}$ .

②  $p$  convex on  $\mathbf{B} \rightarrow \nabla^2 p(\mathbf{x}) \succeq 0$  for all  $\mathbf{x} \in \mathbf{B}$ ,

$$\iff \mathbf{u}^T \nabla^2 p(\mathbf{x}) \mathbf{u} \geq 0, \forall (\mathbf{x}, \mathbf{u}) \in \mathbf{B} \times \mathbf{U},$$

where  $\mathbf{U} := \{\mathbf{u} : \|\mathbf{u}\|^2 \leq 1\}$ .

Hence we have the two "Positivity constraints"

$$\begin{aligned} f(\mathbf{x}) - p(\mathbf{x}) &\geq 0, \quad \forall \mathbf{x} \in \mathbf{B} \\ \mathbf{u}^T \nabla^2 p(\mathbf{x}) \mathbf{u} &\geq 0, \quad \forall (\mathbf{x}, \mathbf{u}) \in \mathbf{B} \times \mathbf{U}. \end{aligned}$$

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$$\iff \mathbf{u}^T \nabla^2 p(\mathbf{x}) \mathbf{u} \geq 0, \forall (\mathbf{x}, \mathbf{u}) \in \mathbf{B} \times \mathbf{U},$$

where  $\mathbf{U} := \{\mathbf{u} : \|\mathbf{u}\|^2 \leq 1\}$ .

Hence we have the two "Positivity constraints"

$$\begin{aligned} f(\mathbf{x}) - p(\mathbf{x}) &\geq 0, \quad \forall \mathbf{x} \in \mathbf{B} \\ \mathbf{u}^T \nabla^2 p(\mathbf{x}) \mathbf{u} &\geq 0, \quad \forall (\mathbf{x}, \mathbf{u}) \in \mathbf{B} \times \mathbf{U}. \end{aligned}$$

## I: Characterizing convex polynomial underestimators

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## II: Characterizing "tightness"

One possibility is to evaluate the  $L_1$ -norm  $\int_{\mathbf{B}} |f(\mathbf{x}) - p(\mathbf{x})| d\mathbf{x}$

$$\rightarrow \int_{\mathbf{B}} (f(\mathbf{x}) - p(\mathbf{x})) d\mathbf{x} = \underbrace{\int_{\mathbf{B}} f(\mathbf{x}) d\mathbf{x}}_{\text{constant}} - \underbrace{\int_{\mathbf{B}} p(\mathbf{x}) d\mathbf{x}}_{\text{linear in } p!}$$

Indeed, writing  $p(\mathbf{x}) = \sum_{\alpha \in \mathbb{N}^n} p_{\alpha} \mathbf{x}^{\alpha}$ ,

$$\int_{\mathbf{B}} p(\mathbf{x}) d\mathbf{x} = \sum_{\alpha \in \mathbb{N}^n} p_{\alpha} \underbrace{\int_{\mathbf{B}} \mathbf{x}^{\alpha} d\mathbf{x}}_{\gamma_{\alpha}},$$

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where  $\gamma_{\alpha}$  is known (and easy to compute)!

Hence computing the **best degree- $d$  convex polynomial underestimator** of  $f$  reduces to solve the **CONVEX** optimization problem:

$$\begin{aligned} \mathbf{P} : \quad \rho &= \inf_{p \in \mathbb{R}[\mathbf{x}]_d} \sum_{\alpha \in \mathbb{N}_d^n} p_\alpha \gamma_\alpha \\ \text{s.t.} \quad & f(\mathbf{x}) - p(\mathbf{x}) \geq 0, \forall \mathbf{x} \in \mathbf{B} \\ & \mathbf{u}^T \nabla^2 p(\mathbf{x}) \mathbf{u} \geq 0, \forall (\mathbf{x}, \mathbf{u}) \in \mathbf{B} \times \mathbf{U}. \end{aligned}$$



which has an optimal solution  $p^* \in \mathbb{R}[\mathbf{x}]_d$

Hence we replace  $\mathbf{P}$  with the **HIERARCHY of SEMIDEFINITE PROGRAMS**

$$\rho_\ell = \inf_{\mathbf{p} \in \mathbb{R}[\mathbf{x}]_d} \sum_{\alpha \in \mathbb{N}_d^n} p_\alpha \gamma_\alpha$$

$$\text{s.t. } f - \mathbf{p} = \sigma_0(\mathbf{x}) + \sum_{j=1}^m \sigma_j(\mathbf{x}) g_j(\mathbf{x}) + \sigma_{m+1}(M - \|\mathbf{x}\|^2)$$

$$\mathbf{u}^T \nabla^2 \mathbf{p}(\mathbf{x}) \mathbf{u} = \psi_0(\mathbf{x}, \mathbf{u}) + \sum_{j=1}^m \psi_j(\mathbf{x}, \mathbf{u}) g_j(\mathbf{x})$$

$$+ \psi_{m+1}(\mathbf{x}, \mathbf{u})(M - \|\mathbf{x}\|^2) + \psi_{m+2}(\mathbf{x}, \mathbf{u})(1 - \|\mathbf{u}\|^2);$$

$$\deg(\sigma_j, \psi_j) \leq 2\ell,$$

parametrized by  $\ell$ .

→ and optimal solution  $p_\ell^* \in \mathbb{R}[\mathbf{x}]_d$ .

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## Theorem (Lass & T. Phan Thanh (JOGO 2013))

$\rho_\ell \rightarrow \rho$  and  $p_\ell^* \rightarrow p^* \in \mathbb{R}[\mathbf{x}]_d$ , as  $\ell \rightarrow \infty$

→ Provides the best results in the comparison:

Guzman, Y. A; Hasan, M. M. F.; Floudas, C. A: *Computational Comparison of Convex Underestimators for Use in a Branch-and-Bound Global Optimization Framework*, Optimization in Science and Engineering; Springer, 2014; pp 229-246.

# III. Applications in probability

Let  $\mathbf{K} \subseteq \mathbb{R}^n$ ,  $\mathbf{S} \subset \mathbf{K}$  be Borel subsets, and  $\Gamma \subset \mathbb{N}^n$ .

Finding an **upper bound** (if possible **optimal**) on **Prob** ( $\mathbf{X} \in \mathbf{S}$ ), the probability that a  $\mathbf{K}$ -valued random variable  $\mathbf{X} \in \mathbf{S}$ , given some of its moments  $\gamma = \{\gamma_\alpha\}$ ,  $\alpha \in \Gamma \subset \mathbb{N}^n$  ....

.... is equivalent to solving:

$$\sup_{\mu \in M(\mathbf{K})} \{ \mu(\mathbf{S}) \mid \int_{\mathbf{K}} X^\alpha d\mu = \gamma_\alpha, \quad \alpha \in \Gamma \}$$

- $M(\mathbf{K})$  is the (convex) set of **probability measures** on  $\mathbf{K} \subseteq \mathbb{R}^n$ .
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Notice that writing

$$\mu = \nu + \varphi, \quad \text{with}$$

- $\varphi$  supported on  $S$ , and
- $\nu$  supported on  $K \setminus S$ ,

equivalently, one has to solve

$$\mathbf{P} : \sup_{\nu, \varphi} \left\{ \int_S 1 d\varphi \mid \int_{K \setminus S} X^\alpha d\nu + \int_S X^\alpha d\varphi = \gamma_\alpha, \quad \alpha \in \Gamma \right\}$$

(since we maximize  $\varphi(S)$ , one may take  $\nu$  supported on  $K$ .)

Assume that  $\Gamma \subset \mathbb{N}_d^n$ . Then the dual of  $\mathbf{P}$  reads:

$$\mathbf{P}^* : \inf_{p_\alpha} \left\{ \sum_{\alpha \in \Gamma} p_\alpha \gamma_\alpha : p \geq 1 \text{ on } S; p \geq 0 \text{ on } K \right\}$$

where  $p \in \mathbb{R}[\mathbf{x}]_d$  is a polynomial

$$\mathbf{x} \mapsto p(\mathbf{x}) = \sum_{\alpha \in \mathbb{N}_d^n} p_\alpha \mathbf{x}^\alpha; \quad p_\alpha = 0 \quad \forall \alpha \in \mathbb{N}_d^n \setminus \Gamma.$$

Let  $K$  and  $S \subset K$  be compact semi-algebraic sets:

$$\begin{aligned} K &= \{ \mathbf{x} \in \mathbb{R}^n : h_k(\mathbf{x}) \geq 0, \quad k = 1, \dots, p \} \\ S &= \{ \mathbf{x} \in \mathbb{R}^n : g_j(\mathbf{x}) \geq 0, \quad j = 1, \dots, m \} \end{aligned}$$

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Again the basic idea of the moment-SOS approach is to replace the positivity constraints

$$p - 1 \geq 0 \text{ on } S; \quad p \geq 0 \text{ on } K$$

with the SOS-positivity certificates

$$p(\mathbf{x}) - 1 = \sigma_0(\mathbf{x}) + \sum_{k=1}^p \sigma_k(\mathbf{x}) h_k(\mathbf{x})$$
$$p(\mathbf{x}) = \psi_0(\mathbf{x}) + \sum_{j=1}^m \psi_j(\mathbf{x}) g_j(\mathbf{x})$$

for some SOS polynomial  $(\sigma_k) \subset \mathbb{R}[\mathbf{x}]_t$  and  $(\psi_j) \subset \mathbb{R}[\mathbf{x}]_t$ ,



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☞ One ends up in solving the hierarchy of semidefinite programs indexed by  $t \in \mathbb{N}$ :

$$\rho_t = \min_{p \in \mathbb{R}[\mathbf{x}]_d} \sum_{\alpha \in \Gamma} p_\alpha \gamma_\alpha : \quad \text{subject to:}$$

$$p_\alpha = 0, \quad \forall \alpha \in \mathbb{N}_d^n \setminus \Gamma$$

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$$\sigma_k, \psi_j \text{ SOS}; \quad \deg(\sigma_k h_k) \leq 2t; \quad \deg(\psi_j g_j) \leq 2t;$$



# The volume of a basic-semi algebraic set

Let  $S \subset \mathbb{R}^n$  be a compact basic semi-algebraic set. Let  $K$  be a BOX  $[0, a]^n$  containing  $S$  and let:

$$\gamma_\alpha = \int_K X^\alpha dx = \frac{a^{n+|\alpha|}}{\prod_{k=1}^n (1 + \alpha_k)!}, \quad \forall \alpha \in \mathbb{N}^n$$

## Theorem

*The (Lebesgue) volume of the set  $S$  is obtained as:*

$$\sup_{\nu, \varphi} \left\{ \int_S 1 d\varphi \mid \int_S X^\alpha d\varphi + \int_K X^\alpha d\nu = \gamma_\alpha, \quad \alpha \in \mathbb{N}^n \right\}$$

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## PROBLEM 4: Measures with given marginals

Let  $\mathbf{K}_j \subset \mathbb{R}^{n_j}$ ,  $j = 1, \dots, p$ , and  $\mathbf{K} := \mathbf{K}_1 \times \mathbf{K}_2 \cdots \times \mathbf{K}_p \subset \mathbb{R}^n$ , and with natural projections  $\pi_j : \mathbf{K} \rightarrow \mathbf{K}_j$ ,  $j = 1, \dots, p$ .

Let  $\nu_j$  be a given Borel measure on  $\mathbf{K}_j$ ,  $j = 1, \dots, p$

For a measure  $\mu$  on  $\mathbf{K}$ , denote  $\pi_j \mu$  its **marginal** on  $\mathbf{K}_j$ , i.e.

$$\pi_j \mu(B) := \mu(\pi_j^{-1}(B)) = \mu(\{x \in \mathbf{K} : \pi_j x \in B\}), \quad B \in \mathcal{B}(\mathbf{K}_j)$$

Consider the optimization problem:

$$f^* = \inf_{\mu \in M(\mathbf{K})} \left\{ \int_{\mathbf{K}} f d\mu \mid \pi_j \mu = \nu_j, \quad j = 1, \dots, p \right\}$$

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with  $M(\mathbf{K})$  being the set of finite Borel measures on  $\mathbf{K}$ .

Generalization of the famous **Monge-Kantorovich transportation** problem, with many other interesting applications, particularly in **Probability**.

If  $\mathbf{K}_j$  is **compact** then the constraint on **marginals**

$$\pi_j \mu = \nu_j$$

is equivalent to the **countably many linear equalities**

$$\int_{\mathbf{K}} X^\alpha d\mu = \int_{\mathbf{K}_j} X^\alpha d\nu_j, \quad \forall \alpha \in \mathbb{N}^{n_j}$$

between **moments** of  $\mu$  and  $\nu_j$  ...

because the space of **polynomials** is **dense** (for the sup-norm) in the space  $C(\mathbf{K}_j)$  of continuous functions on  $\mathbf{K}_j$ .

Hence for each  $d \in \mathbb{N}$ , we may consider the truncated version

$$\rho_d = \inf_{\mu \in M(\mathbf{K})} \left\{ \int_{\mathbf{K}} f d\mu : \int_{\mathbf{K}} X^\alpha d\mu = \underbrace{\int_{\mathbf{K}_j} X^\alpha d\nu_j}_{=\gamma_{j\alpha}}, \alpha \in \mathbb{N}_d^{n_j}; j = 1, \dots, p \right\}$$

with dual

$$\begin{aligned} \rho_d^* &= \sup_{p_j \in \mathbb{R}[X_j]_d} \left\{ \underbrace{\sum_{j=1}^p \int_{\mathbf{K}_j} p_j d\nu_j}_{=: \sum_{\alpha \in \mathbb{N}_d^{n_j}} p_{j\alpha} \gamma_{j\alpha}} : f(X) - \sum_{j=1}^p p_j(X_j) \geq 0 \text{ on } \mathbf{K} \right\} \\ &= \sum_{\alpha \in \mathbb{N}_d^{n_j}} p_{j\alpha} \gamma_{j\alpha} \end{aligned}$$

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Then if  $\mathbf{K} \subset \mathbb{R}^n$  is compact and in the form

$$\mathbf{K} = \{\mathbf{x} \in \mathbb{R}^n : g_j(\mathbf{X}) \geq 0, \quad j = 1, \dots, m\},$$

one replaces the positivity-on- $\mathbf{K}$  constraint

$$f(\mathbf{X}) - \sum_{j=1}^p p_j(\mathbf{X}_j) \geq 0 \quad \text{on } \mathbf{K}$$

with the SOS-based Putinar's certificate

$$f(\mathbf{X}) - \sum_{j=1}^p p_j(\mathbf{X}_j) = \sigma_0(\mathbf{X}) + \sum_{j=1}^m \sigma_j(\mathbf{X}) g_j(\mathbf{X}),$$

with SOS-weights  $(\sigma_j) \subset \mathbb{R}[\mathbf{X}]$ .



... and solve the hierarchy of semidefinite programs

$$\theta_d = \sup_{p_j \in \mathbb{R}[X_j]_d} \left\{ \sum_{j=1}^p \sum_{\alpha \in \mathbb{N}_d^{n_j}} p_{j\alpha} \gamma_{j\alpha} : \right.$$
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☞ One may prove that  $\rho_d \rightarrow f^*$  as  $d \rightarrow \infty$ .

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## IV. Gaussian measures of semi-algebraic sets

Let  $\mu$  be the Gaussian measure on  $\mathbb{R}^n$  with density  $\mathbf{x} \mapsto \exp(-\|\mathbf{x}\|^2)$  and let  $\mathbf{K} \subset \mathbb{R}^n$  be the non necessarily compact basic semi-algebraic set

$$\mathbf{K} = \{\mathbf{x} \in \mathbb{R}^n : g_j(\mathbf{x}) \geq 0, \quad j = 1, \dots, m\}.$$

Goal:

Approximate  $\mu(\mathbf{K})$  as closely as desired

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## Theorem (Lass 2015)

Let  $f \in \mathbb{R}[\mathbf{x}]$  be strictly positive  $\mu$ -a.e. on  $\mathbf{K}$ , and let  $M(\mathbf{K})$  (resp.  $M(\mathbb{R}^n)$ ) be the space of finite Borel measures on  $\mathbf{K}$  (resp.  $\mathbb{R}^n$ ). Then the optimization problem:

$$f_1^* = \sup_{\nu, \phi} \left\{ \int_{\mathbf{K}} f d\phi : \phi + \nu = \mu; \phi \in M(\mathbf{K}), \nu \in M(\mathbb{R}^n) \right\},$$

has a **unique optimal solution**  $(\phi^*, \nu^*) = (\mu_{\mathbf{K}}, \mu - \mu_{\mathbf{K}})$  where  $\mu_{\mathbf{K}}$  is the restriction of  $\mu$  to  $\mathbf{K}$ , that is:

$$\phi^*(B) = \mu_{\mathbf{K}}(B) = \mu(\mathbf{K} \cap B), \quad \forall B \in \mathcal{B}(\mathbb{R}^n).$$

In particular,  $\phi^*(\mathbf{K}) = \mu(\mathbf{K})$ , and  $f^* = \mu(\mathbf{K})$  if  $f = 1$ .

From  $\phi + \nu = \mu$  one deduces  $\phi \leq \mu$  and therefore

$$f^* \leq \int_{\mathbf{K}} f d\mu = \int f d\mu_{\mathbf{K}}.$$

On the other hand the pair  $(\phi^*, \nu^*) = (\mu_{\mathbf{K}}, \mu - \mu_{\mathbf{K}})$  is a feasible solution with associated cost

$$\int_{\mathbf{K}} f d\phi^* = \int f d\mu_{\mathbf{K}},$$

which proves the optimality of  $(\phi^*, \nu^*)$ .

**Uniqueness** is more delicate. Assume there is another optimal solution  $(\phi, \nu)$ . From  $\phi \leq \mu$  one deduces  $\phi \ll \mu$  and so by Radon-Nykodim

$$\phi(B \cap \mathbf{K}) = \int_{B \cap \mathbf{K}} g d\mu \leq \int_{B \cap \mathbf{K}} d\mu, \quad \forall B \in \mathcal{B}(\mathbb{R}^n),$$

for some nonnegative measurable function  $g$ . Hence  $g \leq 1$ ,  $\mu$ -a.e. on  $\mathbf{K}$ .

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On the other hand, by optimality of  $\phi^*$  and  $\phi$ ,

$$\begin{aligned} f^* &= \int_{\mathbf{K}} f d\mu = \int f d\phi^* = \int f d\phi \\ &= \int_{\mathbf{K}} f g d\mu \end{aligned}$$

which implies

$$0 = \int_{\mathbf{K}} f(1 - g) d\mu,$$

Combining this with  $f > 0$  and  $g \leq 1$   $\mu$ -a.e. on  $\mathbf{K}$ , yields  $g = 1$ ,  $\mu$ -a.e. on  $\mathbf{K}$ .

☞ This yields the desired result that  $\phi = \phi^*$ .  $\square$



# A dual view

A possible dual for the above LP is the LP:

$$\rho^* = \inf_{p \in \mathbb{R}[\mathbf{x}]} \left\{ \int_{\mathbf{K}} p d\mu : p \geq f \text{ on } \mathbf{K}; p \geq 0 \text{ on } \mathbb{R}^n \right\},$$

Indeed it trivially holds that  $\rho^* \geq f^*$ .

A tractable version is obtained by replacing:

- the "hard" positivity constraint  $p - f \geq 0$  on  $\mathbf{K}$ , with the positivity-on- $\mathbf{K}$  certificate

$$p - f = \sigma_0 + \sum_{j=1}^m \sigma_j g_j; \quad \sigma_j \text{ is SOS for all } j\}.$$

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so as to obtain the hierarchy of semidefinite approximations

indexed by  $d \in \mathbb{N}$ :

$$\rho_d^* = \inf_{p \in \mathbb{R}[\mathbf{x}]_d} \left\{ \int_{\mathbb{R}^n} p d\mu : p - f = \sigma_0 + \sum_{j=1}^m \sigma_j g_j; \quad p, \sigma_j \text{ all SOS} \right\}$$

where the degree of the SOS  $p, \sigma_j$  is bounded by  $2d$ .

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For every  $d \in \mathbb{N}$ ,  $\rho_d^* \geq f^*$  and  $\rho_d^* \rightarrow f^*$  as  $d \rightarrow \infty$ .

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One may do the same for the complement  $\mathbf{K}^c := \mathbb{R}^n \setminus \mathbf{K}$  as soon as one can write

$$\mathbf{K}^c = \bigcup_{i=1}^p \Omega_i; \quad \mu(\Omega_i \cap \Omega_j) = 0 \quad \forall (i, j)$$

so that  $\mu(\mathbf{K}^c) = \sum_{i=1}^p \mu(\Omega_i)$ . In doing so one obtains for each  $i = 1, \dots, p$  a sequence  $(\theta_{id})_{d \in \mathbb{N}}$  such that

$$\sum_{i=1}^p \theta_{id} \geq \mu(\mathbf{K}^c) \quad \text{and} \quad \lim_{d \rightarrow \infty} \sum_{i=1}^p \theta_{id} = \mu(\mathbf{K}^c) = \mu(\mathbb{R}^n) - \mu(\mathbf{K}).$$

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With  $f = 1$  one obtains  $\underbrace{\mu(\mathbb{R}^n) - \sum_{j=1}^p \theta_{jd}}_{\omega_d^*} \leq \mu(\mathbf{K}) \leq \rho_d^*$  for all  $d$ ,

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# Examples

Let  $n = 2$ , and  $d\mu = \exp(-\|\mathbf{x}\|^2/\sigma) d\mathbf{x}$  and let  $\mathbf{K}$  be the non-convex quadratic

$$\mathbf{x} \mapsto \mathbf{x}^T \mathbf{A} \mathbf{x} = 0.56 x_1^2 + 0.96 x_1 x_2 - 1.24 x_2^2.$$

$$\mathbf{K} = \{(x, y) : (\mathbf{x} - \mathbf{u})^T \mathbf{A} (\mathbf{x} - \mathbf{u}) \leq 1\} \quad (\text{non-compact}),$$

with  $\mathbf{u} = (0.1, 0.5)$  and  $(0.5, 0.1)$ .

$\mathbf{u} = (0.5, 0.1)$			
$\sigma$	$\rho_9^*$	$\omega_9^*$	$100(\rho_9^* - \omega_9^*)/\omega_9^*$
1	2.829605	2.824718	0.17%
0.8	1.876731	1.876609	0.006%
$\mathbf{u} = (0.1, 0.5)$			
$\sigma$	$\rho_9^*$	$\omega_9^*$	$100(\rho_9^* - \omega_9^*)/\omega_9^*$
1	2.989832	2.986599	0.10%
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More details and (non-compact) examples in  
[arXiv:1508.06132](https://arxiv.org/abs/1508.06132).

## Conclusion

- Provides a sequence of **converging upper and lower bounds** on  $\mu(\mathbf{K})$  for **non necessarily compact basic semi-algebraic sets  $\mathbf{K}$** .
- A general methodology not **set- $\mathbf{K}$** -dependent.
- Also works for the exponential measure on the positive orthant  $\mathbb{R}_+^n$ , and in fact any measure  $\mu$  provided that it satisfies Carleman's condition and one knows all its moments.



... but of course ...

With rough basic implementation and present state-of-the-art SDP solvers, one can obtain a few upper and lower bounds only and for dimension  $n = 2$  or  $n = 3$ . For  $d \geq 15$  numerical accuracy problems show up.

-  Some non-trivial tricks (based on Stokes' formula) permit to improve the quality of bounds.
-  Much remains to be done for a better implementation



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# V. Lebesgue decomposition in action

Given two measures  $\mu$  and  $\nu$  on  $\mathbb{R}^n$ ,

one would like to approximate the Lebesgue decomposition

$$\phi + \psi = \mu; \quad \phi \ll \nu; \quad \psi \perp \nu,$$

of  $\mu$  with respect to  $\nu$ .

... based on the sole knowledge of the moments

$$y_\alpha = \int_{\mathbb{R}^n} \mathbf{x}^\alpha d\mu, \quad z_\alpha = \int_{\mathbb{R}^n} \mathbf{x}^\alpha d\nu, \quad \alpha \in \mathbb{N}^n.$$

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of  $\mu$  and  $\nu$ .

By definition of  $\phi$  and  $\psi$ :

$\phi$  has a **DENSITY** w.r.t.  $\nu$  in  $L_1(\nu)$  (called the **Radon-Nikodym derivative** of  $\mu$  w.r.t.  $\nu$ ). That is, there exists a nonnegative measurable function  $f \in L_1(\nu)$  such that:

$$\phi(A) = \int_A f(\mathbf{x}) d\nu(\mathbf{x}), \quad \forall A \in \mathcal{B}(\mathbb{R}^n).$$

## CLAIM: If one assumes that :

- $f$  is in  $L_\infty(\nu)$  (instead of  $L_1(\nu)$ ), and  $\|f\|_\infty < M$  for some  $M$ ,
- Both moment sequences  $(y_\alpha)$  and  $(z_\alpha)$ ,  $\alpha \in \mathbb{N}^n$  satisfy Carleman's condition:

$$+\infty = \sum_{k=1}^{\infty} \left( \int X_i^{2k} d\mu \right)^{-1/2k} = \sum_{k=1}^{\infty} \left( \int X_i^{2k} d\nu \right)^{-1/2k}$$

for all  $i = 1, \dots, n$ .

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# A hierarchy of semidefinite approximations

Denote the moments of  $\mu$  and  $\nu$  by:

$$\mu_\alpha = \int \mathbf{x}^\alpha d\mu, \quad \nu_\alpha = \int \mathbf{x}^\alpha d\nu, \quad \alpha \in \mathbb{N}^n.$$

Let  $\gamma > 0$  be fixed, and consider the **hierarchy of semidefinite programs**  $\mathbf{P}_d$  indexed by  $d \in \mathbb{N}$ :

$$\begin{aligned} \mathbf{P}_d : \quad \rho_d = & \sup_{y, u, v} y_0 \\ \text{s.t.} \quad & y_\alpha + u_\alpha = \mu_\alpha, \quad \forall \alpha \in \mathbb{N}_d^n \\ & y_\alpha + v_\alpha = \gamma \nu_\alpha, \quad \forall \alpha \in \mathbb{N}_d^n \end{aligned}$$

$$\mathbf{M}_d(y), \mathbf{M}_d(u), \mathbf{M}_d(v) \succeq 0$$

Let  $\phi^*$  and  $\psi^*$  be the Lebesgue decomposition of  $\mu$  w.r.t.  $\nu$ , and let  $f^* \in L_1(\nu)$  be the density of  $\phi^*$  w.r.t.  $\nu$ .

### Theorem (Lasserre 2015)

(i) For each  $d \in \mathbb{N}$ , the semidefinite program has an optimal solution  $(y^d, u^d, v^d)$ .

(ii) Moreover as  $d \rightarrow \infty$ , the triplet of sequences  $(y^d, u^d, v^d)$  converges to some triplet of sequences  $(y^*, u^*, v^*)$  where

$$y_\alpha^* = \int \mathbf{x}^\alpha (\gamma \wedge f^*) d\nu = \int \mathbf{x}^\alpha f_\gamma^* d\nu, \quad \forall \alpha \in \mathbb{N}^n.$$

with  $\|f_\gamma^*\|_\infty \leq \gamma$ .

(iii) So if  $f^* \in L_\infty(\nu)$  with  $\|f^*\|_\infty \leq \gamma$ , then

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Let  $n = 2$ ,  $p \in (0, 1)$  and

- $\nu$  is the Gaussian with density  $\mathbf{x} \mapsto \exp(-\|\mathbf{x}\|^2)$ ,
- $\theta$  is the measure uniform distributed on the circle  $\{\mathbf{x} : x_1^2 + x_2^2 = 1\}$ .

Define the measure  $\mu$  to be

$$\mu = p\nu + (1 - p)\theta,$$

so that the Lebesgue decomposition of  $\mu$  w.r.t.  $\nu$  is

$$(\phi, \psi) = (p\nu, (1 - p)\theta).$$



The table below show relative error between the approximate moments  $\mathbf{u} = (u_\alpha)$  of degree 2 and 4, of the singular part  $\psi$  and those of  $p\theta$  computed with moments up to order  $2d = 14$ .

approx. moments	$x_1^2$	$x_1^4$	$x_1^2 x_2^2$	$L_{\mathbf{u}^d}((x_1^2 + x_2^2 - 1)^2)$
p=0.1	0.19%	0.52%	0.53%	0.001
p=0.2	3.7%	8.12%	12.14%	0.16

Same thing but now with  $\nu$  being uniformly supported on the unit box  $[-1, 1]^n$ .

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